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APPLICATION NO.	FILING DATE	FIRST NAMED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.
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09/863,419

05/24/2001

William D. Norcott

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7590

11/20/2008

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EXAMINER

SAEED, USMAAN

ART UNIT

PAPER NUMBER

2166

MAIL DATE

DELIVERY MODE

11/20/2008

PAPER

Please find below and/or attached an Office communication concerning this application or proceeding.

The time period for reply, if any, is set in the attached communication.

Office Action Summary	Application No.	Applicant(s)	
	09/863,419	NORCOTT, WILLIAM D.	
	Examiner	Art Unit	
	USMAAN SAEED	2166	

-- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --

Period for Reply

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE 3 MONTH(S) OR THIRTY (30) DAYS, WHICHEVER IS LONGER, FROM THE MAILING DATE OF THIS COMMUNICATION.

- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133). Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

Status

- 1) ☒ Responsive to communication(s) filed on 08 August 2008.
- 2a) ☒ This action is **FINAL**. 2b) ☐ This action is non-final.
- 3) ☐ Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

Disposition of Claims

- 4) ☒ Claim(s) 1 and 3-10 is/are pending in the application.
- 4a) Of the above claim(s) _____ is/are withdrawn from consideration.
- 5) ☐ Claim(s) _____ is/are allowed.
- 6) ☒ Claim(s) 1 and 3-10 is/are rejected.
- 7) ☐ Claim(s) _____ is/are objected to.
- 8) ☐ Claim(s) _____ are subject to restriction and/or election requirement.

Application Papers

- 9) ☐ The specification is objected to by the Examiner.
- 10) ☒ The drawing(s) filed on 24 May 2001 is/are: a) ☒ accepted or b) ☐ objected to by the Examiner.
Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).
Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).
- 11) ☐ The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

Priority under 35 U.S.C. § 119

- 12) ☐ Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).
- a) ☐ All b) ☐ Some * c) ☐ None of:
1. ☐ Certified copies of the priority documents have been received.
 2. ☐ Certified copies of the priority documents have been received in Application No. _____.
 3. ☐ Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).

* See the attached detailed Office action for a list of the certified copies not received.

Attachment(s)

- | | |
|--|---|
| 1) <input checked="" type="checkbox"/> Notice of References Cited (PTO-892) | 4) <input type="checkbox"/> Interview Summary (PTO-413) |
| 2) <input type="checkbox"/> Notice of Draftsperson's Patent Drawing Review (PTO-948) | Paper No(s)/Mail Date. _____ |
| 3) <input type="checkbox"/> Information Disclosure Statement(s) (PTO/SB/08) | 5) <input type="checkbox"/> Notice of Informal Patent Application |
| Paper No(s)/Mail Date _____ | 6) <input type="checkbox"/> Other: _____ |

DETAILED ACTION

Response to Amendment

1. Receipt of Applicant's Amendment, filed 08/08/2008 is acknowledged.
Claims 1, 6, 7, and 8 have been amended.

Claim Rejections - 35 USC § 103

2. The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:

(a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negated by the manner in which the invention was made.

This application currently names joint inventors. In considering patentability of the claims under 35 U.S.C. 103(a), the examiner presumes that the subject matter of the various claims was commonly owned at the time any inventions covered therein were made absent any evidence to the contrary. Applicant is advised of the obligation under 37 CFR 1.56 to point out the inventor and invention dates of each claim that was not commonly owned at the time a later invention was made in order for the examiner to consider the applicability of 35 U.S.C. 103(c) and potential 35 U.S.C. 102(e), (f) or (g) prior art under 35 U.S.C. 103(a).

Claims 1, 3-4 and 6 are rejected under 35 U.S.C. 103(a) as being unpatentable over **William D. Norcott** ('**Norcott**' hereinafter), USP, 5,848,405 in view of **David W. Moore**. ('**Moore**' hereinafter), (USP, 6,594,676).

With respect to claim 1,

Norcott teaches **a method for change data capture** (see col. 1, lines 64 to col. 2, line 1), **comprising the steps of:**

executing a database statement (see col. 6, lines 54-55) **to copying a recovery log** (new data for refresh processing purposes, the server process deletes the ROWID range from the ROWID range table. Updates a redo log to indicate the changes made to the range table ensures that identification of the new data can be recovered in the event of database crash, see col. 6, lines 23-31 and col. 5, lines 59-61, **Norcott**),

that contains change data for all transactions performed from a source object in a first system to a object in a second system (updating data in data warehouses identifies new data by storing a plurality of new data records into the database at contiguous storage locations, and storing range data that specifies the range of the contiguous storage locations. Hence, new data is identified in a database by storing the start location and the end location of a contiguous range of data locations, enabling server processes to access the new data records based on the stored range data without the necessity of accessing another database table on a row by row basis, See col. 2, Lines 1-8, and figure 6, **Norcott**),

storing the recovery log that contains the change data (storing a plurality of new data records into the database at contiguous storage locations, See col. 1, Lines 65-66, and figure 6, **Norcott**). The redo log/recovery log 620 is also being updates with the changes.

obtaining at least some the of the change data from the recovery log based in part on a column in the recovery log that holds data representing when a transaction has been committed (The redo log 620 records changes made by all transactions within the database. Since the range data, including the start and end ROWID values, are database objects, updating the redo log 620 to indicate the changes made to the range table ensures that identification of the new data 610b can be recovered in the event of a database crash that affects the data in the ROWID range table 612, see Col 5, lines 54-62, **Norcott**. The first table entry specifies an extent having data records with contiguous storage location identifiers having a range within the low ROWID equal to A up to and including the high ROWID equal to C. The second, third, and fourth table entries specify extents 202 having data records at the contiguous ranges of ROWIDs D-F, G-J, and K-M, respectively. If desired, the table column includes a flag ("&") appended to the table name to signify that the next row entry specifies the location of the next extent storing the new data, see Col 6 lines 64-67 and Col 7, Lines 1-6, **Norcott**),

inserting the at least some change data into a plurality of change tables in the second system in a single transaction, where each change table includes a column corresponding to the column in the recovery log (updating data in data warehouses identifies new data by storing a plurality of new data records into the database at contiguous storage locations, and storing range data that specifies the range of the contiguous storage locations. Hence, new data is identified in a database by storing the start location and the end location of a contiguous range of data locations,

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enabling server processes to access the new data records based on the stored range data without the necessity of accessing another database table on a row by row basis, See col. 2, Lines 1-8, and figure 6, **Norcott**),

Norcott does not explicitly indicate the claimed “**copying**,” “**obtaining data from the columns in recovery log**” and “**inserting change data into plurality of change tables**”.

However, **Moore** discloses “**copying**,” “**obtaining data from the columns in recovery log**” and “**inserting change data into plurality of change tables**” (as the database system 200 comprises one or more CADS 214 designated CADS1 to CADSn to indicate a variance in the number of CADS 214 in the system 200. The CADS 214 contains records reflecting change data from one or more logs 204 for a certain span of time. A single CADS 214 may further reflect updates for one or more databases 206 See col. 8, lines 47-55 and figures 2 and 3, **Moore**).

These lines and figures teach that CADS utility 306 reads log data 204 to create one or more CADS 214. Examiner interprets CADS containing change data from recovery log as a result of copying from the recovery log. A single CADS is further being used to insert updates to plurality of databases 206. Moore further provides a single read of a CADS having data for more than one database data set.

It would have been obvious to one of ordinary skill in the art at the time the invention was made to combine the teaching of the cited references because **Moore’s** teaching would have allowed **Norcott** to provide a recovery utility apparatus for expediting recovery time during failure of one or more database data sets by eliminating

the need to sequentially read each backup copy and CADS for each CADS associated with a database requiring recovery.

Claim 6 is essentially the same as claim 1, except it recites the claimed invention as a computer readable medium bearing instructions and is rejected for the same reasons as applied hereinabove.

As to claim 3,

Norcott teaches **renaming a column heading for a source column of the source object and providing the renamed column heading for the source column as a column as a column heading in a change column of the database object** (all database records belonging to the table names "sales" and having ROWIDs in the range between x and y (inclusive) are identified as new records having just been inserted into the database (e.g., table portion 610b) see, Col 6, Lines 44-48, **Norcott**). (The first table entry specifies an extent having data records with contiguous storage location identifiers having a range within the low ROWID equal to A up to and including the high ROWID equal to C. The second, third, and fourth table entries specify extents 202 having data records at the contiguous ranges of ROWIDs D-F, G-J, and K-M, respectively. If desired, the table column includes a flag ("&") appended to the table name to signify that the next row entry specifies the location of the next extent storing the new data, see Col 6 lines 64-67 and Col 7, Lines 1-6, **Norcott**),

As to claim 4,

Norcott teaches **generating the database statement to store the change data in the database object** (If the new data records are stored entirely within a single group of data records having a contiguous sequence of ROWIDs, then the summary refresh process is completed after the server process deletes the ROWID range from the ROWID range table, see col. 6, lines 27-32 and lines 50-58, **Norcott**).

3. Claims 5, 7-10 are rejected under 35 U.S.C. 103(a) as being unpatentable over **William D. Norcott** ('**Norcott**' hereinafter), USP, 5,848,405 in view of **David W. Moore**. ('**Moore**' hereinafter), (USP, 6,594,676) as applied to claims 1, 3-4 and 6 further in view of **Rauer et al.** ('**Rauer**' hereinafter), (USP, 6,161,103).

As to claim 5,

Norcott teaches **shipping change data from the recovery log of an on-line transaction processing (OLTP) system** (The source of the data is an online transaction processing (OLTP) database and OLTP databases provide a mechanism for exporting [shipping] data from the database into a static file, see col. 4, lines 20-25, **Norcott**).

Norcott does not explicitly indicate the claimed "**staging system**".

Rauer discloses claimed staging system (SQL statements are issued to the source system and the results are loaded into the staging tables. The staging tables had

been created as a result of block. Once the staging tables have been loaded, the data can then be moved into the datamart, see col. 10, lines 49-53, **Rauer**).

It would have been obvious to one ordinary skill in the data processing art, at the time of the present invention, to combine the teachings of the cited references because the staging system of **Rauer's** teachings would have allowed to Norcott and Moore to create databases, loading and accessing data in the databases as suggested by Ruer, see col. 1, lines 46-47. Further, staging system as taught by Rauer improves to perform set of commands to execute the creation of the aggregate tables (see col. 3, lines 19-20, Rauer).

With respect to claim 7,

Norcott teaches **a method of change data capture** (see col. 1, lines 64 to col. 2, line 1), **comprising the steps of:**

shipping change data for at least one transaction that has been performed on an on-line transaction processing (OLTP) system from a recovery log that contains change data for all transactions performed to an on-line transaction processing (OLTP) system (Updates a redo log to indicate the changes made to the range table ensures that identification of the new data can be recovered in the event of database crash, see col. 5, lines 59-61, **Norcott**), (The redo log 620 records changes made by all transactions within the database, See col. 5, Lines 56-57, **Norcott**,) and (The source of the data is an online transaction processing (OLTP) database and OLTP

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databases provide a mechanism for exporting [shipping] data from the database into a static file, see col. 4, lines 20-25, **Norcott**); and

Norcott teaches **performing the steps** (see col. 5, lines 59-61) of: **copying the recovery log from the on-line transaction processing (OLTP) system** (new data for refresh processing purposes, the server process deletes [extract] the ROWID range from the ROWID range table. Updates a redo log to indicate the changes made to the range table ensures that identification of the new data can be recovered in the event of database crash, see col. 6, lines 29-31 and col. 5, lines 59-61, **Norcott**); and

storing the recovery log in a database object (the start and end ROWID values are database objects and stored in step 406, see col. 5, lines 58-66, Fig. 4, **Norcott**),

database object having at least one control column (The first table entry specifies an extent having data records with contiguous storage location identifiers having a range within the low ROWID equal to A up to and including the high ROWID equal to C. The second, third, and fourth table entries specify extents 202 having data records at the contiguous ranges of ROWIDs D-F, G-J, and K-M, respectively. If desired, the table column includes a flag ("&") appended to the table name to signify that the next row entry specifies the location of the next extent storing the new data, see Col 6 lines 64-67 and Col 7, Lines 1-6, **Norcott**),

obtaining change data from the recovery log stored in the first object based in part on a column in the recovery log that holds data representing when a transaction has been committed (The redo log 620 records changes made by all

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transactions within the database. Since the range data, including the start and end ROWID values, are database objects, updating the redo log 620 to indicate the changes made to the range table ensures that identification of the new data 610b can be recovered in the event of a database crash that affects the data in the ROWID range table 612, see Col 5, lines 54-62, **Norcott**. The first table entry specifies an extent having data records with contiguous storage location identifiers having a range within the low ROWID equal to A up to and including the high ROWID equal to C. The second, third, and fourth table entries specify extents 202 having data records at the contiguous ranges of ROWIDs D-F, G-J, and K-M, respectively. If desired, the table column includes a flag ("&") appended to the table name to signify that the next row entry specifies the location of the next extent storing the new data, see Col 6 lines 64-67 and Col 7, Lines 1-6, **Norcott**),

inserting the change data into a plurality of change tables in the second database, where the change data is inserted into each of the plurality of change tables in a single transaction, where each change table includes a column corresponding to the column in the recovery log (updating data in data warehouses identifies new data by storing a plurality of new data records into the database at contiguous storage locations, and storing range data that specifies the range of the contiguous storage locations. Hence, new data is identified in a database by storing the start location and the end location of a contiguous range of data locations, enabling server processes to access the new data records based on the stored range data

without the necessity of accessing another database table on a row by row basis, See col. 2, Lines 1-8, and figure 6, **Norcott**),

said change data representing modifications that has been performed to a plurality of source tables of the on-line transaction processing (OLTP) system and that to the change tables in the second database (see col. 1, Lines 64-67 to col. 2, Lines 1-8 and Figure 6, **Norcott**).

Norcott does not explicitly indicate the claimed “**staging system,**” “**copying,**” “**obtaining data from the columns in recovery log**” and “**inserting change data into plurality of change tables**”.

However, **Moore** discloses “**copying,**” “**obtaining data from the columns in recovery log**” and “**inserting change data into plurality of change tables**” (as the database system 200 comprises one or more CADS 214 designated CADS1 to CADSn to indicate a variance in the number of CADS 214 in the system 200. The CADS 214 contains records reflecting change data from one or more logs 204 for a certain span of time. A single CADS 214 may further reflect updates for one or more databases 206 See col. 8, lines 47-55 and figures 2 and 3, **Moore**).

These lines and figures teach that CADS utility 306 reads log data 204 to create one or more CADS 214. Examiner interprets CADS containing change data from recovery log as a result of copying from the recovery log. A single CADS is further being used to insert updates to plurality of databases 206. Moore further provides a single read of a CADS having data for more than one database data set.

It would have been obvious to one of ordinary skill in the art at the time the invention was made to combine the teaching of the cited references because **Moore's** teaching would have allowed **Norcott** to provide a recovery utility apparatus for expediting recovery time during failure of one or more database data sets by eliminating the need to sequentially read each backup copy and CADS for each CADS associated with a database requiring recovery.

Norcott and Moore teach elements of claim 7 as noted above but do not explicitly teach a “**staging system.**”

Rauer discloses claimed staging system (SQL statements are issued to the source system and the results are loaded into the staging tables. The staging tables had been created as a result of block. Once the staging tables have been loaded, the data can then be moved into the datamart, see col. 10, lines 49-53, **Rauer**).

It would have been obvious to one ordinary skill in the data processing art, at the time of the present invention, to combine the teachings of the cited references because the staging system of **Rauer's** teachings would have allowed to Norcott and Moore to create databases, loading and accessing data in the databases as suggested by Ruer, see col. 1, lines 46-47. Further, staging system as taught by Rauer improves to perform set of commands to execute the creation of the aggregate tables (see col. 3, lines 19-20, Rauer).

With respect to claim 8,

Norcott teaches a **method of change data capture** (see col. 1, lines 64 to col. 2, line 1), **comprising the steps of:**

shipping change data for at least one transaction that has been performed on an on-line transaction processing (OLTP) system from a recovery log that contains change data for all transactions performed to the on-line transaction processing (OLTP) system (Updates a redo log to indicate the changes made to the range table ensures that identification of the new data can be recovered in the event of database crash, see col. 5, lines 59-61, **Norcott**), (The redo log 620 records changes made by all transactions within the database, See col. 5, Lines 56-57, **Norcott**), and (The source of the data is an online transaction processing (OLTP) database and OLTP databases provide a mechanism for exporting [shipping] data from the database into a static file, see col. 4, lines 20-25, **Norcott**); and

Norcott teaches **performing** (see col. 1, lines 6-8) **the steps of: registering the recovery log with a log viewer** (Updates a redo log to indicate the changes made to the range table ensures that identification of the new data can be recovered in the event of database crash, see col. 5, lines 59-61, **Norcott**);

generating a SQL statement to extract the change data from the recovery log (After all the new records have been added to the "sales" table it is possible to identify the new records by using the ROWID range table, for example by processing the SQL select statement:

SELECT*FROM sales

WHERE (ROWID>=X) AND (ROWID<=Y)

The summary refresh process can access the new data by processing such a select statement after obtaining the values of x and y from the ROWID range table, see col. 6, lines 50-58, **Norcott**); and

executing the SQL statement (see col. 6, lines 50-55, **Norcott**), **thereby copying the recovery log from the on-line transaction processing (OLTP) system** (new data for refresh processing purposes, the server process deletes [extract] the ROWID range from the ROWID range table. Updates a redo log to indicate the changes made to the range table ensures that identification of the new data can be recovered in the event of database crash, see col. 6, lines 29-31 and col. 5, lines 59-61, and figure 6 **Norcott**); and

storing the recovery log in a database object (the start and end ROWID values are database objects and stored in step 406, see col. 5, lines 58-66, Fig. 4, **Norcott**),

database object having at least one control column (The first table entry specifies an extent having data records with contiguous storage location identifiers having a range within the low ROWID equal to A up to and including the high ROWID equal to C. The second, third, and fourth table entries specify extents 202 having data records at the contiguous ranges of ROWIDs D-F, G-J, and K-M, respectively. If desired, the table column includes a flag ("&") appended to the table name to signify that the next row entry specifies the location of the next extent storing the new data, see Col 6 lines 64-67 and Col 7, Lines 1-6, **Norcott**),

obtaining change data from the recovery log stored in the first object based in part on a column in the recovery log that holds data representing when a transaction has been committed (The redo log 620 records changes made by all transactions within the database. Since the range data, including the start and end ROWID values, are database objects, updating the redo log 620 to indicate the changes made to the range table ensures that identification of the new data 610b can be recovered in the event of a database crash that affects the data in the ROWID range table 612, see Col 5, lines 54-62, **Norcott**. The first table entry specifies an extent having data records with contiguous storage location identifiers having a range within the low ROWID equal to A up to and including the high ROWID equal to C. The second, third, and fourth table entries specify extents 202 having data records at the contiguous ranges of ROWIDs D-F, G-J, and K-M, respectively. If desired, the table column includes a flag ("&") appended to the table name to signify that the next row entry specifies the location of the next extent storing the new data, see Col 6 lines 64-67 and Col 7, Lines 1-6, **Norcott**),

inserting the change data into a plurality of change tables in the second database, where the change data is inserted into each of the plurality of change tables in a single transaction, where each change table includes a column corresponding to the column in the recovery log (updating data in data warehouses identifies new data by storing a plurality of new data records into the database at contiguous storage locations, and storing range data that specifies the range of the contiguous storage locations. Hence, new data is identified in a database by storing the

start location and the end location of a contiguous range of data locations, enabling server processes to access the new data records based on the stored range data without the necessity of accessing another database table on a row by row basis, See col. 2, Lines 1-8, and figure 6, **Norcott**),

said change data representing modifications that has been performed to a plurality of source tables of the on-line transaction processing (OLTP) system and that to the change tables in the second database (see col. 1, Lines 64-67 to col. 2, Lines 1-8 and Figure 6, **Norcott**).

Norcott does not explicitly indicate the claimed “**staging system,**” “**copying,**” “**obtaining data from the columns in recovery log**” and “**inserting change data into plurality of change tables**”.

However, **Moore** discloses “**copying,**” “**obtaining data from the columns in recovery log**” and “**inserting change data into plurality of change tables**” (as the database system 200 comprises one or more CADS 214 designated CADS1 to CADSn to indicate a variance in the number of CADS 214 in the system 200. The CADS 214 contains records reflecting change data from one or more logs 204 for a certain span of time. A single CADS 214 may further reflect updates for one or more databases 206 See col. 8, lines 47-55 and figures 2 and 3, **Moore**).

These lines and figures teach that CADS utility 306 reads log data 204 to create one or more CADS 214. Examiner interprets CADS containing change data from recovery log as a result of copying from the recovery log. A single CADS is further

being used to insert updates to plurality of databases 206. Moore further provides a single read of a CADS having data for more than one database data set.

It would have been obvious to one of ordinary skill in the art at the time the invention was made to combine the teaching of the cited references because **Moore's** teaching would have allowed **Norcott** to provide a recovery utility apparatus for expediting recovery time during failure of one or more database data sets by eliminating the need to sequentially read each backup copy and CADS for each CADS associated with a database requiring recovery.

Norcott and Moore teach elements of claim 8 as noted above but do not explicitly teach a “**staging system.**”

Rauer discloses claimed staging system (SQL statements are issued to the source system and the results are loaded into the staging tables. The staging tables had been created as a result of block. Once the staging tables have been loaded, the data can then be moved into the datamart, see col. 10, lines 49-53, **Rauer**).

It would have been obvious to one ordinary skill in the data processing art, at the time of the present invention, to combine the teachings of the cited references because the staging system of **Rauer's** teachings would have allowed to **Norcott and Moore** to create databases, loading and accessing data in the databases as suggested by Ruer, see col. 1, lines 46-47. Further, staging system as taught by Rauer improves to perform set of commands to execute the creation of the aggregate tables (see col. 3, lines 19-20, **Rauer**).

As to claim 9,

Norcott teaches renaming a column heading for a source column of the source object and providing the renamed column heading for the source column as a column as a column heading in a change column of the database object (all database records belonging to the table names "sales" and having ROWIDs in the range between x and y (inclusive) are identified as new records having just been inserted into the database (e.g., table portion 610b) see, Col 6, Lines 44-48, **Norcott**). (The first table entry specifies an extent having data records with contiguous storage location identifiers having a range within the low ROWID equal to A up to and including the high ROWID equal to C. The second, third, and fourth table entries specify extents 202 having data records at the contiguous ranges of ROWIDs D-F, G-J, and K-M, respectively. If desired, the table column includes a flag ("&") appended to the table name to signify that the next row entry specifies the location of the next extent storing the new data, see Col 6 lines 64-67 and Col 7, Lines 1-6, **Norcott**),

As to claim 10,

Norcott teaches wherein the on-line transaction processing (OLTP) system are provided by different database vendors employing a different, incompatible internal implementation (The source of the data is an online transaction processing (OLTP) database and OLTP databases provide a mechanism for exporting [shipping] data from the database into a static file. The static file then loaded by the server process

into the database table and enable the database to processes for update, see col. 4, lines 20- 30, **Norcott**).

Response to Arguments

4. Applicant's arguments have been considered but are moot in view of the new ground(s) of rejection.

In these arguments applicant relies on the amended claims and not the original ones. See above rejections for the arguments.

Claims must be given the broadest reasonable interpretation during examination and limitations appearing in the specification but not recited in the claim are not read into the claim (See M.P.E.P. 2111 [R-I]).

Conclusion

5. Applicant's amendment necessitated the new ground(s) of rejection presented in this Office action. Accordingly, **THIS ACTION IS MADE FINAL**. See MPEP § 706.07(a). Applicant is reminded of the extension of time policy as set forth in 37 CFR 1.136(a).

A shortened statutory period for reply to this final action is set to expire THREE MONTHS from the mailing date of this action. In the event a first reply is filed within TWO MONTHS of the mailing date of this final action and the advisory action is not mailed until after the end of the THREE-MONTH shortened statutory period, then the shortened statutory period will expire on the date the advisory action is mailed, and any

extension fee pursuant to 37 CFR 1.136(a) will be calculated from the mailing date of the advisory action. In no event, however, will the statutory period for reply expire later than SIX MONTHS from the date of this final action.

Contact Information

6. Any inquiry concerning this communication or earlier communications from the examiner should be directed to Usmaan Saeed whose telephone number is (571)272-4046. The examiner can normally be reached on M-F 8-5.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Hosain Alam can be reached on (571)272-3978. The fax phone number for the organization where this application or proceeding is assigned is 571-273-8300.

Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see <http://pair-direct.uspto.gov>. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free).

Usmaan Saeed
Patent Examiner
Art Unit: 2166

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Hosain Alam
Supervisory Patent Examiner

US
November 17, 2008

/Hosain T Alam/

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